Extrinsic faulting in 3C close-packed crystal structures: computational mechanics analysis

Ernesto Estevez-Rams, Raimundo Lora-Serrano, Arbelio Penton-Madrigal and Massimo Nespolo

Extrinsic faulting in 3C close-packed crystal structures: computational mechanics analysis

Ernesto Estevez-Rams, a* Raimundo Lora-Serrano, b Arbelio Penton-Madrigal c and Massimo Nespolo c

a Facultad de Física – Instituto de Ciencia y Tecnología de Materiales (IMRE), Universidad de la Habana, San Lazaro y L. CP 10400 C., Habana, Cuba, b Instituto de Física, Universidade Federal de Uberlândia, Uberlândia 38400-902, Minas Gerais, Brazil, and c Université de Lorraine, CNRS, CRM2, Nancy, France. *Correspondence e-mail: estevez@fisica.uh.cu

Extrinsic faulting has been discussed previously within the so-called difference method and random walk calculation. In this contribution it is revisited under the framework of computational mechanics, which allows expressions to be derived for the statistical complexity, entropy density and excess entropy as a function of faulting probability. The approach allows one to compare the disordering process of an extrinsic fault with other faulting types. The \( \varepsilon \)-machine description of the faulting mechanics is presented. Several useful analytical expressions such as probability of consecutive symbols in the Hägg coding are presented, as well as hexagonality. The analytical expression for the pairwise correlation function of the layers is derived and compared with results previously reported. The effect of faulting on the interference function is discussed in relation to the diffraction pattern.

1. Introduction

Close-packed structures (CPS) are OD (order–disorder) structures built by stacking hexagonal layers in the direction perpendicular to the layer (Đurović, 1997). The stacking ambiguity arising from the two possible positions of a layer with respect to the previous one leads to a theoretically infinite number of possible polytypes if no constraint is made on the periodicity. However, by far the commonest ones are the cubic close-packed or 3C (Bravais lattice of type \( cF \)) and the hexagonal close-packed or 2H (Bravais lattice of type \( hP \)). These are MDO (maximum degree of order) polytypes, meaning that the structure contains the minimal number of layer triples, quadruples etc. (one, in both cases).

CPS usually exhibit some kind of planar disorder or stacking faults, which, when viewed as a disruption in the otherwise periodic arrangement of layers, can be analysed as non-interacting defects. This is the basis of the so-called random faulting model (RFM), which has been the most widely used model of faulting in layer structures, dating back to the early times in diffraction analysis (Wilson, 1942; Wagner, 1957; Warren, 1969). The idea of the RFM is to consider certain types of faulting, such as intrinsic (removal of a layer from the sequence), extrinsic (addition of a layer in the sequence) and twinning (change of orientation in the sequence), assigning to each a fixed probability of occurrence, independent from the density of faulting and neglecting any spatial interaction between the faults present in the material. This simplifying assumption means that the RFM, if suitable, should be at a very low density of faulting; in this case, during the derivation of the corresponding expression for the
displacement correlation function between layers (also known as the pairwise correlation) and subsequently the diffraction equation, it is justified to drop all terms above linear in the faulting probabilities. Analytical expressions are then found; for such mathematical development readers are referred to the classical works of Warren (1969) for deformation and twin faulting in f.c.c. (face-centred cubic) and h.c.p. (hexagonal close-packed) structures.

Besides the assumption of a low density of defects, the RFM also assumes that faults, when they occur, go through the whole coherently diffracting domain, avoiding the need to account for the appearance of partial dislocations. Further, faulting is considered to happen along the stacking direction but not along any direction that is crystallographically equivalent, as in any cubic crystal, but this is no longer the case if faulting occurs in one of the four. The reader is referred to Estevez-Rams et al. (2007) for further historical account of the subject.

In recent years there have been attempts to extend the mathematical applicability of the model, without modifying its fundamental assumptions. First, Velterop et al. (2000) have observed that, even for a low density of faulting, the assumption of only one faulting direction is an unrealistically simplifying assessment of the diffraction behaviour, which can lead to misleading conclusions. Another issue is the need to accommodate larger, more realistic, faulting densities within the model. Even if the physical assumption of non-interacting faulting is too strong for a larger probability of faulting, it is still interesting to pursue such an extension for several reasons. The RFM can be used as a reference model for other approaches. The fact that only one parameter for each faulting type is needed makes it very attractive in the practical analysis of materials. Additionally, the RFM can be used as a suitable starting model in computer simulations of faulting. In this case, the need for a good starting proposal is essential in the convergence and convergence speed of numerical calculations.

In the last few years, independently, Varn et al. (Varn & Canright, 2001; Varn, 2001; Varn et al., 2002; Varn & Crutchfield, 2004) and Estevez-Rams et al. (2008) have attempted to rewrite the RFM in a modern framework, using a hidden Markov model (HMM) description of the faulting dynamics. The more ambitious idea is to go beyond the faulting model and try to understand the disordering process in layer structures as a dynamical process of a system capable of performing (physical) computation and, in this sense, able to store and process information (Crutchfield, 1992). The attractiveness of the proposal is that such an approach can make use of a powerful set of tools, developed within the study of complexity, grounded in information theory concepts such as Shannon entropy and mutual information. This framework is known as computational mechanics and has been used in a wide range of subjects (Crutchfield, 2012).

A first attempt to use the HMM description of random faulting was made by Estevez-Rams et al. (2008) for intrinsic and twinning faults. Their analysis allowed calculation of the displacement correlation function and the diffraction equation for the whole range of faulting probabilities. They also derived a useful expression concerning the hexagonality of the stacking event, the average size of cubic and hexagonal neighbourhood blocks, the correlation length, all as a function of the faulting probabilities. While correct, this approach is ad hoc in nature, only applicable to the problem considered by them, namely using as starting structure the 3C layer ordering, and working through the appropriate equations.

A more recent breakthrough came from the work of Riechers et al. (2015), who proved that calculation of the pairwise correlation function could be systematized in an elegant way, allowing its applicability to a wide number of situations such as those found in close-packed arrangements.

The result of Riechers et al. (2015) opens the possibility to study, in a systematic way, the RFM for different types of faulting and their combinations, something which proved to be at least cumbersome and, in certain instances, intractable using previous tools. This is what we intend to do in this contribution for extrinsic faults. Extrinsic faulting has been dealt with before (Johnson, 1963; Warren, 1963; Lele et al., 1967; Holloway & Klamkin, 1969; Howard, 1977; Howard & Kuwano, 1979; Takahashi, 1978).

The main goal of the article is to report several analytical expressions for disorder of extrinsic faulted CPS. These expressions relate disorder magnitudes, such as those derived within mechanical computation, with the extrinsic faulting probability which in turn allows comparison with similar expressions derived for twin and deformation faults already reported (Estevez-Rams et al., 2008). Also, a closed analytical expression for the probability of finding different stacking sequences in the faulted structure is reported and from there an expression is derived for the hexagonality and the average length of perfectly coherent f.c.c. sequences within the CPS. The analytical expression of the pair correlation function as a function of faulting probability is derived and its decaying and oscillation behaviour are discussed. Finally, the expression for the interference function is reported and peak shift and asymmetry as a result of the extrinsic faulting are discussed.

First, the main concepts used and the notation are introduced.

2. Order and disorder in close-packed stacking arrangements and the pairwise correlation function

In the OD structures built from hexagonal layers, the layers can be found only in three positions perpendicular to the stacking direction, which are usually labelled A, B and C.
Close-packed is the constraint that two layers which bear the same letter, and are thus exactly overlapped in the projection along the stacking direction, cannot occur consecutively. According to this description, the ideal f.c.c. structure is described by $ABCABCAB \ldots$ sequences (Verma & Krishna, 1966), while the ideal h.c.p. structure has a stacking order described by $ABABABA \ldots$ and the double hexagonal close-packed (d.h.c.p.) structure in turn is described by the stacking $ABCBABCB \ldots$.

An equivalent, and less redundant, coding is the Hägg code (Hagg, 1943), where a pair of consecutive layers is given a plus (or 1) symbol if they form a ‘forward’ sequence $AB$, $BC$ or $CA$, and a minus (or 0) sign otherwise.\(^1\) There is a one-to-one relation between both codings (Estevez-Rams et al., 2005). It is also important to introduce a three-layer hexagonal environment as one where a layer $X$ has the two adjacent layers in the same position (e.g. $ABA$, $ACA$, $BAB$, $BCB$, $CAB$, $CBC$); if a layer environment is not hexagonal then it is cubic. A hexagonal environment is denoted by a letter $h$ and a cubic environment by a letter $k$; this is the basis of the Jagodzinski coding of the stacking arrangement and, as before, there is a one-to-one correspondence between the $ABC$ coding and the Jagodzinski coding (Estevez-Rams & Martinez-Mojicar, 2008). Hexagonality then refers to the fraction of hexagonal environments in the stacking sequence. Also, it can be easily checked that when in the Hägg code the pair of characters 10 or 01 is found, a hexagonal environment is found.

Faulting is generically meant as a disruption of the ideal periodic ordering of a stacking arrangement and therefore constitutes a defect in the structure. In CPS, the simplest types of faults that are usually considered are (i) deformation faults, which are jogs in the otherwise perfect periodic sequence, (ii) extrinsic or double deformation fault, which is the insertion of an extraneous layer in the perfect sequence and (iii) twin faults, which cause reversions in the stacking ordering. In what follows, the probability of the occurrence of a deformation fault will be denoted by $a$, of an extrinsic fault by $\gamma$, while the probability of the occurrence of a twin faulting will be denoted by $\beta$.

The pair correlation function between layers, known as the pairwise correlation function $Q_2(\Delta)$, is the key to calculating the effect of the stacking arrangement on the diffraction intensity (Estevez-Rams et al., 2001, 2003). Consider a stacking direction and sense, $Q_2(\Delta)$, where $\xi = [c, a, s]$ is the probability of finding two layers, $\Delta$ layers apart, the first displaced with respect to the second one as (i) $\xi = c: A - B$ or $B - C$ or $C - A$; (ii) $\xi = a: B - A$ or $C - B$ or $A - C$; and (iii) $\xi = s: A - A$ or $B - B$ or $C - C$.\(^2\) It should be noted that $Q_2(1) = 0$ due to the close-packed constraint and $Q_2(0) = 1$ by construction.

It is possible, for any of the described codings $ABC$, Hägg and Jagodzinski, to construct an HMM describing a broad range of both ordered and disordered stacking processes. An HMM description comprises a finite, or at least enumerable, set of states $\mathcal{S}$ and the associated initial set of the probability $\pi_\alpha$ of being in each state; the set of transition matrices $T$; and a set of symbols drawn from a finite alphabet $\mathcal{A}$. Each transition matrix $T^{ij}$ is a square matrix with the number of rows equal to the number of states, where each entry $T^{ij}_{k}$ represents the probability of jumping from state $i$ to state $j$, while emitting the symbol $v \in \mathcal{A}$. The HMM transition matrix $T$ is defined as the sum of the $T^{ij}$ over all symbols $v$ in the alphabet. Fig. 1 shows the HMM for the f.c.c., h.c.p. and d.h.c.p. stacking structures. For further details the reader is referred to previous papers on the subject (Varn & Canright, 2001; Varn, 2001; Varn et al., 2002; Varn & Crutchfield, 2004; Estevez-Rams et al., 2008).

When seen through an HMM description, stacking arrangements are cast as an information processing system that sequentially outputs symbols as it makes transitions between states. Care must be taken not to necessarily ascribe physical meaning to the states in the HMM description; they are, in general, only meaningful in the sense of the information processing structure of the system. The system output is an infinite string of characters $Y = \ldots v_{-2}v_{-1}v_{0}v_{1}v_{2} \ldots$, each character $v_i \in \mathcal{A}$. For the purpose of analysis it is common, at a given point, to divide the output string into two halves, the left half $Y = \ldots v_{-2}v_{-1}$ is known as the past, while the right half $Y = v_{0}v_{1}v_{2} \ldots$ is called the future\(^3\) (Varn & Canright, 2001; Varn, 2001; Varn et al., 2002; Varn & Crutchfield, 2004).

There can be many HMMs describing the same process. The minimal HMM describing the system dynamics is considered to be optimal in the sense of using fewer resources while providing the best predictive power and will be the one relevant in this contribution; such a model is called an $\varepsilon$-machine (Crutchfield, 1992, 2012). The $\varepsilon$-machine has the important property, among others, of unifiability which means that, from a given state, the emitted symbols determine unambiguously the transition to another state.

---

\(^1\) An alternative notation by Nabarro-Frank (Frank, 1951) uses $\gamma$ and $\Delta$ for + (or 1) and $-$ (or 0) respectively.

\(^2\) The notation is that used by Varn et al. (Varn & Canright, 2001; Varn et al., 2002), where $c$ stands for ‘cyclic’, $a$ stands for ‘anti-cyclic’ and $s$ for ‘same’.

---

**Figure 1**

The HMM in the Hägg notation for the perfect f.c.c. structure $\ldots 11111 \ldots$; the h.c.p. structure $\ldots 1010101 \ldots$; and d.h.c.p. structure $\ldots 1100110011 \ldots$. Labels at the transition arcs are of the form $s|p$ where $s$ is the emitted symbol and $p$ the probability of making the transition from one state to the other.
Let us denote, following the common use of bras and kets in physics, by $\langle \pi |$ the vector of state probabilities and by $| 1 \rangle$ a vector of $1$s. If the HMM description is known, then the probability of any finite sequence $\nu^N = u_i u_{i+1} u_{i+2} \ldots u_{i+N-1}$ will be given by

$$P(\nu^N) = \langle \pi | T^{[u_1]} T^{[u_2]} \ldots T^{[u_{N-1}]} | 1 \rangle,$$

(1)

where in this case $\langle x| A | y \rangle$ is a real number resulting from the scalar product between vectors and matrices.

Several information theory magnitudes can be defined once the minimal HMM description of the stacking process is known. Shannon defined information entropy $H(X)$ for an event set $X$ with a discrete probabilities distribution $p(X)$ as (Arndt, 2001)

$$H(X) = - \sum_i p(X) \log p(X),$$

(2)

where the sum is taken over all the probability distribution, and here and in what follows the logarithm is taken base two which makes the units of entropy bits.

For the $\varepsilon$-machine, the statistical complexity $C_\mu$ is defined as the Shannon entropy over the HMM states,

$$C_\mu = H(S) = - \sum_i p_i \log p_i,$$

(3)

where $p_i$ is the stationary probability of the $i$th state in the minimal HMM description and the sum is over all state probabilities. $C_\mu$ measures the amount of information the system stores.

Excess entropy $E$ is also used to characterize the system information processing capabilities and is used as a measure of predictability, defined as the mutual information between the left half and the right half in the system output,

$$E = H(\hat{\Upsilon}) + H(\hat{\Upsilon}) - H(\Upsilon).$$

(4)

Entropy density $h_\mu$ (Arndt, 2001) is defined as

$$h_\mu = \lim_{N \rightarrow \infty} \frac{H(\Upsilon^N)}{N},$$

(5)

when such limits exist, with $\Upsilon^N$ denoting all substrings of $\Upsilon$ of length $N$. $h_\mu$ is used to answer how random the process is (Crutchfield & Feldman, 2003).

Finally, Riechers et al. (2015) described a procedure for computing the pairwise correlation function from the transition matrices, which can be summarized as follows:

(i) The HMM of the stacking process in the ABC notation is given together with $\{A, S, \pi_0, T\}$. If this description is given in the Hägg coding then the expansion to the $ABC$ coding must be performed (Riechers et al., 2015).

(ii) The stationary probabilities $\pi$ over the HMM states are calculated as the normalized left eigenvector of the transition matrix $T$ with eigenvalue unity,

$$\langle \pi | = \langle \pi | T.$$

(6)

(iii) The pairwise correlation function follows from the definition and the use of equation (1):

$$Q_\Delta(\Delta) = \sum_{i,j \in A} \langle \pi | T^{[x_i]} T^{-1} T^{[x_j]} | 1 \rangle$$

(7)

where $\hat{\Delta} = \{\hat{c}, \hat{a}, \hat{s}\}$ is a family of permutation functions given by

$$\hat{e}(A) = B \quad \hat{c}(B) = C \quad \hat{a}(C) = A$$

$$\hat{a}(A) = C \quad \hat{a}(B) = A \quad \hat{a}(C) = B$$

$$\hat{s}(A) = A \quad \hat{s}(B) = B \quad \hat{s}(C) = C$$

(8)

and $1$ represents a vector of $1$s [see also equations (20) and (24) in Riechers et al. (2015) for alternative expressions for equation (7)].

3. Extrinsic fault in the face-centred cubic stacking order

An extrinsic fault (also known as a double deformation fault in the case of the f.c.c. structure) in a $3C$ stacking is depicted in Fig. 2 along with a perfect sequence for comparison. It can be seen that in the Hägg code the extrinsic fault is equivalent to the flip (bitwise negation) of two consecutive characters. The probability of occurrence of such faulting will be denoted by $\gamma$. It will be assumed that $\gamma$ can take any value between 0 and 1. Building from the effect of the extrinsic fault over the Hägg code, the HMM of the faulting process is shown in Fig. 3, where it is assumed that the ideal $3C$ structure goes in the $A \rightarrow B \rightarrow C$ sequence. The $p$ state represents the non-faulted condition; as long as the system stays in that state, the output symbol $\nu = 1$ will correspond to the perfect $3C$ structure. If faulting occurs, a $0$ is emitted and the system goes to the $e$ state.

Figure 3

Finite-state automata (HMM) over the Hägg code for the extrinsic fault in the $3C$ stacking. For $\gamma = 0$ and $\gamma = 1$ the HMM describes a perfect $3C$ stacking. The output of the HMM will be of the perfect $3C$ structure as long as it stays in the $p$ state; a transition to the $e$ state signals the occurrence of an extrinsic fault described by, at least, a pair of $0$'s. Observe that the perfect sequence is described by the cyclic $ABC$ sequence corresponding to $1$'s in the Hägg notation; an equivalent HMM could be described for the anti-cyclic $CBA$ sequence by a trivial relabelling of the output symbols.

Figure 2

Extrinsic faulting in the sequence $3C$ along the cubic $[111]$ direction. An ideal periodic sequence is shown for comparison.
state, where a second 0 is printed with certainty while returning to the $p$ state.\(^4\) The HMM of Fig. 3 represents a biased even process [see Appendix A of Crutchfield et al. (2016) and Example D in Varn et al. (2013)].

It should be observed that any sequence with an odd number of 0's cannot be the result of such an HMM. Such a sequence will be called forbidden; moreover, the forbidden sequences are called irreducible as they do not contain a proper subsequence which itself is forbidden. The number of irreducible forbidden sequences in the even process is infinite; in such a case, the process is called a sofic system (Crutchfield & Feldman, 2003). The fact that any sequence from the HMM of the even process contains an even number of 0's has important consequences as will be discussed later on.

The corresponding transition matrix will be given by

\[
T^{[1]} = \begin{pmatrix} \overline{p} & 0 \\ 0 & 0 \end{pmatrix}, \quad T^{[0]} = \begin{pmatrix} 0 & \gamma \\ 1 & 0 \end{pmatrix}, \quad T = \begin{pmatrix} \overline{p} & \gamma \\ 1 & 0 \end{pmatrix},
\]

(9)

where $\overline{p}$ stands for $1-\gamma$. The stationary probabilities over the recurrent states $p$ and $e$ can be calculated following equation (6) which results in

\[
\langle \pi \rangle = \left\{ \frac{1}{1+\gamma}, \frac{\gamma}{1+\gamma} \right\},
\]

(10)

the first value corresponding to the $p$ state.

Hexagonality in terms of computational mechanics has been analysed in a more general context previously (Varn et al., 2007). Hexagonality can be calculated from the probability of occurrence of 01 or 10 in the Hågg code of the sequence. Both probabilities are equal and, from equation (1), given by

\[
P(01) = \langle \pi \rangle T^{[0]} T^{[1]} | 1 \rangle = \frac{\gamma}{1+\gamma},
\]

(11)

from which the hexagonality is given by $2P(01)$. Hexagonality has a maximum value of $2(3-2\sqrt{2}) \simeq 0.343$ at $\gamma = \sqrt{2} - 1 \simeq 0.414$ (Fig. 4a).

The statistical complexity can be derived from equation (3) using equation (10) and is given by

\[
C_A = \frac{1}{1+\gamma} \left[ \log(1+\gamma) - \gamma \log \frac{\gamma}{1+\gamma} \right].
\]

(12)

The logarithm is taken usually in base two and then the unit of $C_A$ is bits. For an $\epsilon$-machine the entropy density is given by (Crutchfield et al., 2016)

\[
h_\mu = -\sum_{k \in A} P(k) \sum_{x \in A} P(x|k) \log P(x|k),
\]

(13)

where $P(a|b)$ means the probability of $a$ conditioned on $b$. The unit of the entropy density is bits per site. The expression for the entropy density will not be derived explicitly and the reader is referred to Crutchfield et al. (2016); the resulting expression for the entropy density is

\[\text{Figure 4} (a) \text{ Hexagonality, entropy density and (b) excess entropy as a function of the faulting probability. Excess entropy has a discontinuity at } \gamma = 1, \text{ where it drops to zero.}
\]

\[
h_\mu = -\frac{1}{1+\gamma} \left[ \gamma \log \gamma + (1-\gamma) \log(1-\gamma) \right].
\]

(14)

The calculation of the excess entropy is more involved and is explained in detail in Appendix A. The result is

\[
E = \frac{1}{1+\gamma} \left[ \log(1+\gamma) - \gamma \log \frac{\gamma}{1+\gamma} \right],
\]

(15)

which is identical to the statistical complexity.

Fig. 4(b) shows the behaviour of the excess entropy as a function of $\gamma$. Observe that, at $\gamma = 1$, the excess entropy has a discontinuity, as $E$ drops to zero when the finite-state automata description has a topological change to a certain process with only one state and emitting always a 0 symbol. This discontinuity is not seen for the entropy density (Fig. 4a) which has a maximum at $\gamma = 1/2(3 - \sqrt{5}) \simeq 0.382$ with value $h_\mu = 0.6942$ bits per site and then smoothly drops to zero as $\gamma$ approaches 1.

The probability of a chain of 0's of length $n$ is given by

\[
P(0^n) = \begin{cases} 
\gamma^n \left( 1 - \frac{2\gamma^n}{1+\gamma} \right) & n = 2l \\
0 & n = 2l + 1
\end{cases}
\]

(16)

For chains of 1's

\[
P(1^n) = \frac{(1-\gamma)^n}{1+\gamma}.
\]

(17)

---

\(^4\)The described dynamics implicitly assume that an inserted layer cannot follow another inserted layer. This case has been approached by Howard (1977).
can be constructed from the H"{a}gg description and is shown in Fig. 3. The subscripts in the states of the HMM are induced in the HMM of the ABC code for the extrinsic fault in the 3C stacking. The HMM can be derived from that of Fig. 3. The subscripts in the states of the HMM are induced in the HMM of the ABC coding describing the extrinsic fault.

Contrary to the previous assumption that entropy density does not determine uniquely the hexagonality, which comes as no surprise as hexagonal neighbourhoods are the result of faulting events, the higher the hexagonality, which comes as no surprise as hexagonal neighbourhoods are the result of faulting events, which in turn implies larger disorder. It can be seen though that entropy density does not determine uniquely the hexagonality of the stacking arrangement. Contrary to the previous behaviour, hexagonality is uniquely determined by the excess entropy. A maximum value of hexagonality is found for an excess entropy of 0.8724 bits.

3.1. The pairwise correlation function

The HMM for the ABC coding corresponding to subsequences starting with A, B and C. Using the same procedure described for the H"{a}gg HMM, the transition matrices can be written and the stationary probability over the recurrent states calculated for the HMM over the ABC coding:

\[ \langle \pi_{abc} \rangle = \frac{1}{3(1+\gamma)} \{ \gamma, \gamma, 1, 1, 1 \}. \]  

where the order in the states has been taken as \{A, B, C, A, B, C\}. Using equation (7) the pairwise correlation follows:

\[ Q_s(\Delta) \]

\[ = \frac{1}{3} \left( 1 + \left( \frac{|p|}{4} \right)^{\Delta} \left\{ \begin{array}{l} \cos(3\varphi) \sin(\Delta \varphi) \\ \sqrt{3(1+\gamma)} \end{array} \right. \right. \]

\[ + \left( \frac{|q|}{4} \right)^{\Delta} \left\{ \begin{array}{l} \cos(3\varphi) \sin(\Delta \varphi) \\ \sqrt{3(1+\gamma)} \end{array} \right. \right. \]

\[ - \left( \frac{|r|}{4} \right)^{\Delta} \left\{ \begin{array}{l} \cos(3\varphi) \sin(\Delta \varphi) \\ \sqrt{3(1+\gamma)} \end{array} \right. \right. \]

\[ = \frac{1}{3} \left[ 1 + Q_s^{[1]}(\Delta) + Q_s^{[2]}(\Delta) \right]. \tag{20} \]

where

\[ r = |r| \exp(i\varphi_r) = \sqrt{i\sqrt{3}(6\gamma - \gamma^2 - 1) - (1+\gamma)^2}, \]

\[ x = (\gamma - 1)(1 - i\sqrt{3}), \]

\[ p = |p| \exp(\varphi_p) = x + \sqrt{2s}, \]

\[ q = |q| \exp(\varphi_q) = x - \sqrt{2s}. \]

The obtained equation is equivalent to the result given by Holloway & Klamkin (1969), as can be seen by comparing numerical results from equation (20) for \( \Delta = 0, 1, 2, 3 \) and

\[ \langle \pi_{abc} \rangle = \frac{1}{3(1+\gamma)} \{ \gamma, \gamma, 1, 1, 1 \}. \]  

Figure 6

Finite-state automata over the ABC code for the extrinsic fault in the 3C stacking. The HMM can be derived from that of Fig. 3. The subscripts in the states of the HMM are induced by the states with the same label in the H"{a}gg HMM.

Figure 5

Hexagonality as a function of (a) entropy density, (b) excess entropy.

From equations (16) and (17) the average length of blocks of 0’s and 1’s can be calculated:

\[ \langle L_0 \rangle = \sum_{n=1}^\infty n P(0^n) = \frac{4\gamma}{(1-\gamma)^2}, \]

\[ \langle L_1 \rangle = \sum_{n=1}^\infty n P(1^n) = \frac{1 - \gamma}{\gamma^2 - 1 + \gamma}. \tag{18} \]

\[ \langle L_0 \rangle = \langle L_1 \rangle \] at \( \gamma = 0.3623. \]

In Fig. 5 hexagonality as a function of excess entropy and entropy density is shown. Higher the excess entropy is, the higher the hexagonality, which comes as no surprise as hexagonal neighbourhoods are the result of faulting events, which in turn implies larger disorder. It can be seen though that entropy density does not determine uniquely the hexagonality of the stacking arrangement. Contrary to the previous behaviour, hexagonality is uniquely determined by the excess entropy. A maximum value of hexagonality is found for an excess entropy of 0.8724 bits.
When comparing with Holloway & Klamkin (1969) results, it must be noticed that in their notation \( Q_s(\Delta) = P(m) \), \( Q_s(\Delta) = Q(m) \) and \( Q_s(\Delta) = R(m) \).

Those reported in equations (35), (36), (37), (38) in Holloway & Klamkin (1969) (making \( \alpha = 0 \)). In turn these authors have shown that their result reduces to that of Johnson (1963). Holloway & Klamkin do not give a close form of \( Q_s(\Delta) \) for \( \Delta > 3 \).

There are two terms in the expression for \( Q_s(\Delta) \), each with an oscillating and a decaying part. Fig. 7(a) shows the behaviour of both decaying terms with faulting probability. \( p \) and \( q \) have a jump (discontinuity) at the same value of \( \gamma \approx 0.1716 \), where the real part of \( r \) has a minimum and the imaginary part has a jump from a negative value to a positive one. Interestingly, the combined plot of both terms results in two smooth continuous curves. At \( \gamma = 0 \), \( p \) is zero while \( q = 1 \), and the oscillating part of the second term in \( Q_s \) dominates. At \( \gamma = 1 \), both \( p \) and \( q \) have the same value of 1 and the combined effect of both oscillating terms determines the pairwise correlation function. For both cases (\( \gamma = 0 \) and \( \gamma = 1 \)), \( Q_s(\Delta) \) reduces to

\[
Q_s(\Delta) = \frac{1}{3} \left[ 1 + 2 \cos \left( \frac{2\pi}{3} \Delta \right) \right],
\]

describing the correlation function for the perfect 3C stacking.

At \( \gamma_0 \), the oscillating part of both terms in \( Q_s(\Delta) \) becomes equal for all values of \( \Delta \). At \( \gamma = \sqrt{2} - 1 \), where the hexagonality reaches its maximum value, the oscillating part of \( Q_s(\Delta) \) is the prevailing one at large \( \Delta \) values. For small (\( \gamma \approx 0 \)) and large values (\( \gamma \approx 1 \)) it is the oscillating part of \( Q_s(\Delta) \) that determines the underlying stacking sequence.

In any case, the lower curve in Fig. 7(a) determines the faster decaying behaviour of the pairwise correlation function, while the upper curve determines the dominant behaviour at larger \( \Delta \) values. Fig. 7(b) shows the correlation lengths derived from both decaying terms. At large values of \( \Delta \) the \( p \) term is the dominant factor in the pairwise correlation function for values of \( \gamma > \gamma_0 \), while the opposite happens at values below \( \gamma_0 \).

A similar deduction made for \( Q_s(\Delta) \) results in

\[
Q_s(\Delta) = \frac{1}{3} \left[ 1 + \left( \frac{|p|}{4} \right)^\Delta \left[ C_p \cos(\Delta \phi_p) + S_p \sin(\Delta \phi_p) \right] \right. \\
\left. + \left( \frac{|q|}{4} \right)^\Delta \left[ C_q \cos(\Delta \phi_q) + S_q \sin(\Delta \phi_q) \right] \right].
\]

with

\[
C_p = \frac{\sqrt{2} \gamma^2 - 4 \gamma + 1}{|r|} \cos \phi_p + 2 \frac{\sqrt{6}}{|r|} \gamma \sin \phi_p - \frac{1}{2},
\]

\[
S_p = \frac{\sqrt{2} \gamma^2 - 4 \gamma + 1}{|r|} \sin \phi_p - 2 \frac{\sqrt{6}}{|r|} \gamma \cos \phi_p + \frac{\sqrt{3}}{2},
\]

\[
C_q = -\frac{\sqrt{2} \gamma^2 - 4 \gamma + 1}{|r|} \cos \phi_q - 2 \frac{\sqrt{6}}{|r|} \gamma \sin \phi_q + \frac{1}{2},
\]

\[
S_q = -\frac{\sqrt{2} \gamma^2 - 4 \gamma + 1}{|r|} \sin \phi_q + 2 \frac{\sqrt{6}}{|r|} \gamma \cos \phi_q + \frac{\sqrt{3}}{2},
\]

\( Q_s(\Delta) \) follows from the normalization condition.

4. The interference function

The diffraction pattern of an OD structure can be decomposed into two contributions: that of the layer and that of the stacking sequence. The reduced diffracted intensities (i.e. once the necessary corrections are applied: Lorentz, polarization, absorption etc.) can be deconvoluted in terms of these two contributions so that the stacking sequence leaves its fingerprint in the form of an interference function showing a periodic distribution of deconvoluted intensities.

In the case of complex sequences like that of micas, in which adjacent layers can be stacked in six different orientations, the interference function has been called the PID (periodic intensity distribution; Nespolo et al. (1999)). For CPS, the situation is simpler because adjacent layers may take only two relative positions. The consequence of extrinsic faulting on the diffracted intensity is visible in the interference function. The
interference function follows from the use of the expressions for $Q_r$, $Q_s$ and $Q_a$ (Estevez-Rams et al., 2001):

$$I(r^*) = 1 + 2 \sum_{\Delta=1}^{N_c-1} A_\Delta \cos(2\pi \Delta l) + B_\Delta \sin(2\pi \Delta l),$$

(21)

where

$$A_\Delta = \left(1 - \frac{\Delta}{N_c}\right) \left[Q_s(\Delta) + \frac{Q_s(\Delta) + Q_a(\Delta)}{2} \cos\left(\frac{2\pi}{3}(h-k)\right)\right],$$

(22)

$$B_\Delta = \sqrt{3} \left(1 - \frac{\Delta}{N_c}\right) \left[Q_s(\Delta) - Q_a(\Delta)\right].$$

The last case is $h-k = 3n+2$ with $n$ an integer; the coefficients are then

$$A_\Delta = \left(1 - \frac{\Delta}{N_c}\right) \left[Q_s(\Delta) - \frac{Q_s(\Delta) + Q_a(\Delta)}{2}\right],$$

$$B_\Delta = \frac{\sqrt{3}}{2} \left(1 - \frac{\Delta}{N_c}\right) \left[Q_s(\Delta) - Q_a(\Delta)\right].$$

An analytical expression for the interference function can be deduced from the above equations but is too long and cumbersome to be of any particular interest.\footnote{In Richers et al. (2014), a more elegant way to deduce the interference function directly from the HMM is derived and could lead to a more manageable expression, as has been rightly pointed out by an anonymous referee.}

The result has been discussed already by Johnson (1963), Warren (1963), Holloway & Klamkin (1969). With increasing faulting probability $\gamma$, the peak asymmetrically broadens, lowers its intensity and shifts (Fig. 8). For $h-k = 1$ (mod 3) the peak originally at $l = 3n + 1$ ($n \in \mathbb{Z}$) shifts towards lower $l$ values, while the opposite occurs for $h-k = 2$ (mod 3) where the peak originally at $l = 3n - 1$ shifts towards higher $l$ values. Additionally, at high faulting probability an additional peak appears near the so-called twin position. For $h-k = 1$ (mod 3) ($h-k = 2$ mod 3) the twin position is at $l = 3n - 1$ ($l = 3n + 1$), the additional peak appears at lower (larger) values of $l$ and gradually shifts towards the twin peak position as $\gamma$ increases while strengthening its intensity and decreasing its broadening. The behaviour of the original peak and that of the twin one are not symmetrical, that is, they do not behave the same for $\gamma$ and $1 - \gamma$, respectively. The non-symmetric behaviour of the peaks can be explained by the non-symmetrical character of the HMM describing extrinsic broadening (Fig. 3). A similar profile for single crystals and the particular case of $\gamma = 1/2$ has been reported by Varn et al. (2013).

If one observes the interference for $\gamma = 0.333$ (Fig. 8), it is too often the case in the literature that peak deformations...
with geometry such as these are fitted with models involving more than one phase. The fact that such distortions can be the result of a single type of faulting that does not lead to any polytype should be taken as a note of warning against introducing too readily new structures in profile fitting.

In Fig. 9 the peak shift and asymmetry as a function of faulting probability are shown. Asymmetry has been defined as the ratio between the half width at half-maximum (HWHM) for the right side \( (W_r) \), divided by the HWHM for the left side \( (W_l) \); by construction, the asymmetry is equal to 1 for a perfect symmetric peak.

For powder diffraction it must be considered that the components of a family of planes like the \{111\} (where all members of the family are crystallographically equivalent for the unfaulited crystal and share the same interplanar distance) are no longer equivalent when faulting occurs. For example, when indexed with respect to the hexagonal axis the \{111\} includes the following planes: \( (0,0,3), (0,0,1), (1,0,0), (0,1,0), (1,1,1) \); the first two are unaffected by extrinsic faulting, the next three are of the type \( h-k=1(mod 3) \) and the last three of the form \( h-k=2(mod 3) \). Thus, when simulating the faulted powder diffraction profiles, each component of a plane family must be considered individually. Fig. 10 shows the powder peak profile for the \{111\} where the components not affected by faulting have been left out. The reader can compare with the single-crystal profiles of Fig. 8.

5. Conclusions

Stacking disorder can be viewed in a number of cases as a dynamical system capable of storing and processing information. From this point of view, it has been shown that an extrinsic fault in the Hågg code is a sofic system, where predictability in the future is linked to long-range memory in the past for faulting probabilities within \( ]0,1[ \). A sofic system, such as the one considered here, has no description as a finite-range Markov process. This inability to describe such a simple faulting process by a finite-range model is interesting, as it is common in the literature to try to model faulting by this type of finite-range Markov model.\(^7\) In spite of this, the HMM model for extrinsic faulting is simple enough; it just belongs to a different type of processing machinery. This is precisely the underlying idea of computational mechanics which attempts to find the less sophisticated model for a given process by climbing up in a given hierarchy of possible computational machines until such a description is found.

This character has several interesting consequences. First, the excess entropy equals the statistical complexity of the system. Excess entropy is linked with the structured output of the system, while statistical complexity measures memory stored in the system. In consequence, structure is linked to memory, a result not surprising once it is acknowledged that the HMM of the process is equivalent to a biased even process. In an even process, the occurrence of consecutive 0’s has to be tracked completely, to determine in which state the system is. As increasing faulting probability means longer runs of 0’s, excess entropy grows monotonically with increasing \( \gamma \). Excess entropy has a discontinuity at \( \gamma = 1 \), where the topology of the HMM changes to a one-state system with certainty in the output and therefore zero \( E \).

Entropy density, on the other hand, is a smooth function of faulting probability in all the probability range. Entropy density has a maximum at \( \gamma \approx 0.382 \), near the maximum of the hexagonality, but a slightly larger value. Extrinsic faulting, as treated here, implies that no faulting probability changes the underlying periodic sequence: no phase transformation happens. Hexagonality reaches at \( \gamma = \sqrt{2} - 1 \approx 0.414214 \) a maximum of \( 2(3 - 2\sqrt{2}) \approx 0.34314 \) and therefore the system is always more ‘cubic’ than hexagonal.

In the text, several useful analytical expressions have been derived for different entropic magnitudes, probabilities, lengths and correlations, all as a function of the faulting probability \( \gamma \). To the knowledge of the authors, such expressions have not been reported before.

The pairwise correlation function of the layers has been derived and from there the interference function was obtained. The correlation function is composed of two terms, each with a decaying and oscillating part. The numerical values of the obtained expression coincide with those that can be found using previous treatments. The shift and asymmetric broadening of the reflections as a result of extrinsic faulting were also discussed.

APPENDIX A

Calculation of the excess entropy

In order to calculate the excess entropy the mixed-state representation of the system dynamics must be deduced. To understand what the mixed-state representation is, the HMM description must be viewed as an instance of an HMM (Upper,

---

\(^7\) We thank one of the anonymous referees for her/his enlightening comment on this issue.
The state transition matrix will be

\[ W = \begin{pmatrix} 0 & 2\gamma & 1-\gamma & 0 \\ \frac{1}{1+\gamma} & 0 & 2\gamma & 0 \\ 0 & 0 & 1-\gamma & \gamma \\ 0 & 0 & 0 & 1 \end{pmatrix} \]

with eigenvalues

\[ \Lambda = \{1, -\gamma, -\sqrt{\gamma}, \sqrt{\gamma}\} \]

The projection operator \( W_\lambda \), for each eigenvalue, is obtained using

\[ W_\lambda = \prod_{\xi \in \Lambda, \xi \neq \lambda} \frac{W - \xi I}{\lambda - \xi}. \]

\( I \) represents the identity matrix and the product avoids the singularity in the denominator. The results are

\[ W_1 = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{\gamma}{\gamma+1} & \frac{\gamma}{\gamma+1} \\ 0 & 0 & \frac{\gamma}{\gamma+1} & \frac{\gamma}{\gamma+1} \\ 0 & 0 & \frac{\gamma}{\gamma+1} & \frac{\gamma}{\gamma+1} \end{pmatrix} \]

\[ W_{-\gamma} = \begin{pmatrix} \frac{1}{2} & -\sqrt{\gamma} & \frac{1}{2} & \frac{1}{2} \sqrt{\gamma} - \frac{1}{2} \sqrt{\gamma} + \frac{1}{2} \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \]

\[ W_{-\sqrt{\gamma}} = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} & \frac{1}{2} & \frac{1}{2} \sqrt{\gamma} - \frac{1}{2} \sqrt{\gamma} + \frac{1}{2} \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \]

\[ W_{\sqrt{\gamma}} = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} & \frac{1}{2} & \frac{1}{2} \sqrt{\gamma} + \frac{1}{2} \sqrt{\gamma} - \frac{1}{2} \sqrt{\gamma} + \frac{1}{2} \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \]

Defining

\[ \langle \delta \rangle = \{1, 0, 0, 0\} \]

then

\[ |H(W^4)| = - \sum_{\delta \in S} \sum_{x \in [0,1]} \langle \delta \rangle \langle W^{(x)} | 1 \rangle \log \langle \delta \rangle |W^{(x)}| 1 \rangle, \]

and the excess entropy follows from

\[ E = \sum_{\lambda \in \Lambda, |\lambda| < 1} \frac{1}{|\lambda|} \langle \delta_{\lambda \pi_{\text{mix}}} | W_\lambda | H(W^4) \rangle \]

which is equation (8) from Crutchfield et al. (2016).
Acknowledgements

This work was partially financed by FAPEMIG under project BPV-00047-13 and received computational infrastructure support under project APQ-02256-12. EER wishes to thank the Université de Lorraine for a visiting professor grant. He would also like to acknowledge financial support under the PVE/CAPES grant 1149-14-8 which allowed him to visit the UFU. RLS is grateful for the support of CNPq through projects 309647/2012-6 and 304649/2013-9. We would like to thank the anonymous referees for their careful reading and the valuable suggestions that greatly improved the manuscript.

References

Varn, D. P. (2001). PhD thesis, the University of Tennessee, USA.